

The double (raster) IRQ method

Theory

This method is one of the more easy methods to grasp when it comes to stable timing. The following explanation of the method assumed you've accustomed yourself to the previous articles in this series.

As stated earlier a raster IRQ will trigger at cycle 0 on the chosen scanline, but due to the fact that the CPU must finish the current opcode + that the CPU needs 7 cycles to actually save the state and move the PC to the IRQ-handle we'll end up somewhere on cycle 7-14 (assuming no sprites on this raster line).

So, due to the finish-current-opcode problem we have a jitter of 7 cycles. However the shorted opcode (in time) is a NOP f.e., so wouldn't it be nifty if we knew that the CPU was executing NOPs when the raster IRQ occurred? That would be still 7 cycles for IRQ-setup, but only a jitter of 1 cycle. I.e. we would end up on cycle 7 or 8.

A 2 cycle NOPs only gives a 1-cycle jitter because either the CPU has just finished the current opcode and can thus jump immediately to the IRQ-handle, or it is half way through the NOP and must stall and additional one cycle to finish it.

Now, how to guarantee the CPU to execute NOPs when the IRQ occurs?

It's quite easy, let's configure a raster IRQ at raster line X to pull the CPU out of the main-code. This IRQ will jitter with 7 cycles at most since the main code can be anything. Now, immediately let's configure the VIC to trigger another raster IRQ at line X+1. We're almost at X+1 but quite not there yet so let's just fill up with NOPs. This would ensure that when the second IRQ occurs we know 100% that the CPU will be executing NOPs and thus give us a jitter of only 1 cycle at line X+1.

Remember that the C64 has two irq pointer locations? One at \$fffe in the ROM memory which is due to hardware design and one at \$0314 in RAM which is just a memory location the kernal uses to indirect call a user irq handler. We can use both pointer to accomplish the task of irq jitter correction. The first irq, which contains the "nop" list, is called by the \$0314 pointer with a cpu port value of \$37 or \$36. The second irq which just has a jitter of 1 cycle is called by the rom pointer at \$fffe with a cpu port value of \$35. The trick is to set up these pointers beforehand and within the irq code just toggle the cpu port value to toggle the irq targets. The number of nops can be determined by just removing one after each other until the part crashes. Then add again one nop. This is how it is done in the following code example where \$0314 points to "irq0" and \$fffe points to "irq". The three pla just removes the pc and status from the stack which were pushed at the moment the second irq was triggered. However after the second irq ends we want to jumpback to the main program and not somewhere in the middle of the nop list. The inc \$d012 in the first irq must come early so we still in the rasterline previous to the rasterline in which the second irq should be triggered. The same applies for the dec \$d012 in the second irq for similar reasons. After the three pla some timing opcodes follow. These assure that we are at the end of a rasterline when the lda \$d012 is being performed. The compare tests, if the load occurred in rasterline N or N+1 and wastes a cycle in the first case. As an effect the lines after the branch are executed always at the same rasterbeam position. The irq is stable, now.


```
    sta $dc0d
    bit $dc0d
    lda #$81
    sta $d01a
    lda #<irq0
    sta $0314
    lda #>irq0
    sta $0315
    lda #<irq
    sta $fffe
    lda #>irq
    sta $ffff
    lda #$38
    sta $d012
    lda #$36
    sta $01
    lsr $d019
    rts
.ENDPROC
```

The `lsr $d019` at the end of the `init` is performed to clear the `irq` request if for some reason a `rasterirq` request is rised during `init`.

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